# Rank Monotonicity in Undirected Networks

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# Is it good for me to get...

- A follower on Twitter?
- Formally, it is an arc towards me in a directed graph
- A friend on Facebook?
- and spectral measures of centrality
- For example: closeness centrality, eigenvector centrality, PageRank...

• Formally, an edge between me and someone else in an undirected graph

• To make the "good" part formal, in this talk we consider classic geometric

# For Example, With (In)Degree...

- If I get a new follower or a new friend, my score increases
- Moreover, if someone has a smaller score (fewer friends or followers than me), their score remains smaller
- Even more is true, if someone else has the same score as me (i.e., number of followers of friends) they will have a lesser score afterwards
- This properties are obvious for (in)degree
- Can we show that they are still true while we switch to more "sophisticated" scores?

# Score and Rank Monotonicity

- Score monotonicity = something good happens
- Rank monotonicity = nothing bad happens
- than y
- nothing bad happens, something good happens
- Undirected case: must work at both ends (Boldi, Furia & Vigna 2021)

• Score monotonicity: if you add an arc  $x \rightarrow y$ , the score of y increases (Sabidussi 1966)

• Rank monotonicity: if you add an arc  $x \rightarrow y$ , no vertex with a score lower than or equal y can get a score higher than y (Chien, Dwork, Kumar, Simon & Sivakumar 2004)

• Strict rank monotonicity: additionally, vertices with score equal to y get a score lower

# **Geometric Centrality Measures**

- Given a graph, we compute a score using some function of the distances
- Closeness centrality (Bavelas 1948): reciprocal of sum of all distances

- Harmonic centrality (Beauchamp 1965): sum of reciprocals of all distances  $\sum_{y \neq x} \frac{1}{d(y, x)}$
- Betweenness (Anthonisse, 1971; Freeman, 1977)

 $\frac{1}{\sum_{y} d(y, x)}$ 

- the adjacency matrix

- (Actually stated for chess tournaments using  $A\mathbf{1}^{\mathsf{T}}$ )

• Given a graph, we compute a score using some eigenvector associated with

Eigenvector centrality (Landau 1895): just take the dominant eigenvector

• Motivation: If the graph is a voting graph, **1**A is majority voting (indegree)

• We can refine this: let's weight the voters using majority voting:  $(\mathbf{1}A)A = \mathbf{1}A^2$ 

• Or refine again, but  $\mathbf{1}A^k$  oscillates, so Landau proposes to find a positive  $\mathbf{v}$ such that  $\mathbf{v}A = \lambda \mathbf{v}$  (and indeed under mild hypotheses  $\mathbf{1}A^k$  tends to such a  $\mathbf{v}$ )



- Seeley's index (1951): reputation is recursive
- Motivation: kindergarten data about child i liking child j
- Idea: global reputation should be defined recursively:  $s_i = \sum_{j \to i} s_j / d_j$
- This is equivalent to sP = s, where P is A with  $\ell_1$ -normalized rows (i.e., divide by the outdegree)
- Important idea: reputation is divided among people you endorse
- Equivalently, the steady state of the natural Markov chain on the graph (AKA "simplified PageRank"—PageRank without the damping factor)

- Katz's index (1953)
- Follows Landau's idea, but using a summation rather than a limit
- He computes  $\mathbf{1} + \mathbf{1}\alpha A + \mathbf{1}\alpha^2 A^2 + \mathbf{1}\alpha^3$
- One can use a generic border condition v (Hubbell 1965)
- a *right* dominant eigenvector of A (Vigna 2016, using Brauer 1952)

$${}^{3}A^{3} + \cdots = \mathbf{1}\sum_{n\geq 0} \alpha^{n}A^{n} = \mathbf{1}(\mathbf{1} - \alpha A)^{-1}$$

• Note:  $\alpha$  must be less than the inverse of the spectral radius;  $\mathbf{1}A^k$  is also the number of incoming paths of length k, and that is the original formulation

•  $v(1 - \alpha A)^{-1}$  is equal to the dominant eigenvector of  $\alpha A + (1 - \alpha)e^{T}v$ , where e is

- time
- Steady state of a perturbed Markov
- Modulo a factor  $(1 \alpha)$ ,  $v(1 \alpha P)^{-1} =$
- In both cases, when  $\alpha$  tends to its limiting value we go back

PageRank (Page, Brin, Motwani & Winograd 1998); different formulations in

chain: 
$$p = p(\alpha P + (1 - \alpha)1^{T}v)$$

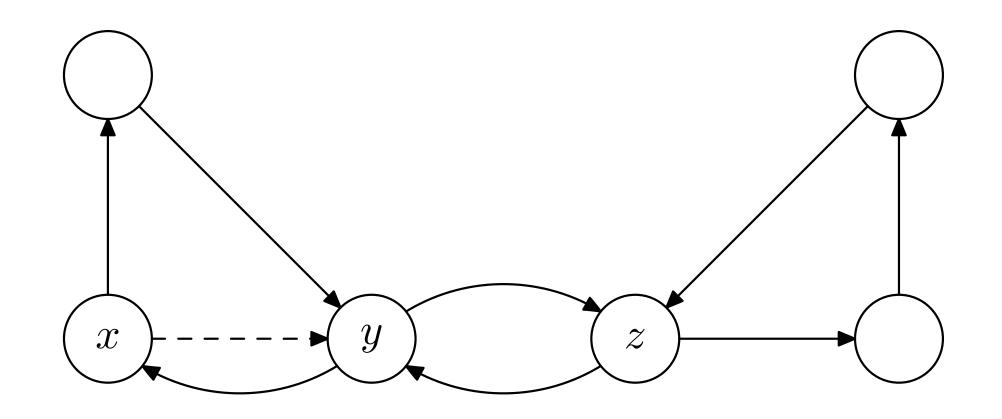
$$\mathbf{v}\sum_{n\geq 0}\alpha^{n}P^{n} = \mathbf{v} + \mathbf{v}\alpha P + \mathbf{v}\alpha^{2}P^{2} + \mathbf{v}\alpha^{3}P^{3} + \cdots$$

So, Seeley's index is to PageRank as eigenvector centrality is to Katz's index

# Getting a New Follower

# The Directed Case: Geometric Centralities

- Not surprisingly: we shorten a path
- More surprisingly: closeness is not strictly rank monotone



One more reason to ditch closeness in favor of harmonic)

On strongly connected graphs, both score monotone and rank monotone

# The Directed Case: Eigenvector Centrality

- (Boldi & Vigna 2004)
- dominant eigenvector increasing more than any other coordinate
- So at least on strongly connected graphs the intuition is correct

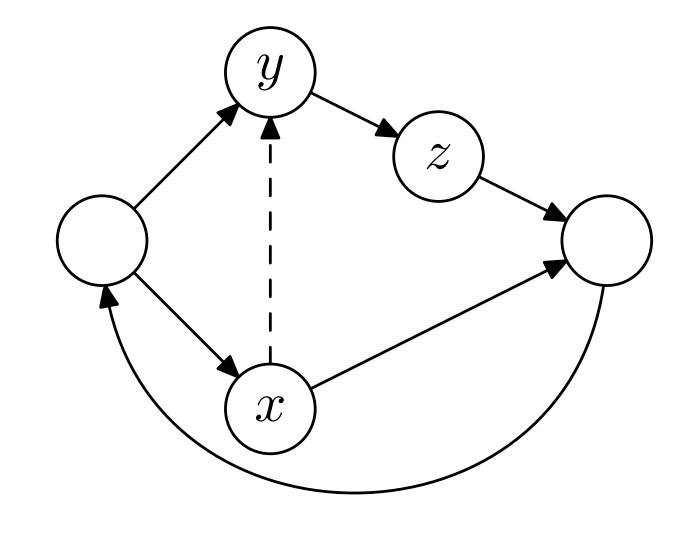
• Eigenvector centrality is strictly rank monotone on strongly connected graphs

 Elsner, Johnson & Neumann 1982: any nonnegative (and overall nonzero) increase in row *i* of a nonnegative matrix results in coordinate *i* of the



# The Directed Case: Seeley's Index

- Score and rank monotone only on strongly connected graphs
- Both results proved only for regular Markov chains in Chien, Dwork, Kumar, Simon & Sivakumar 2004
- First result does not need aperiodicity; second result can be proved using our results for PageRank and taking the limit
- Not strictly rank monotone: y and z maintain the same score



- Score monotonicity and loose rank monotonicity in Chien, Dwork, Kumar, Simon & Sivakumar 2004 under regularity assumptions
- We remove all hypotheses assuming just that the preference vector is positive, and prove strict rank monotonicity
- Or, loose rank monotonicity under the only hypothesis that x has nonzero score
- Our proofs cover a class of spectral centrality measures of the form  $v \sum_{n \ge 0} \alpha^n M^n$
- Same proof for PageRank and Katz's index
- Strategy: use the Sherman–Morrison formula to move the perturbation in the preference vector, and then argue using properties of *M*-matrices

- arc  $x \rightarrow y$
- and 1/(d+1) in position y
- Then the Sherman–Morrison formula expresses  $(1 \alpha P')^{-1}$  as an (ugly)
- Then,  $v(1 \alpha P')^{-1} = (v + c\delta)(1 \alpha P)^{-1}$

Let P and P' be the normalized adjacency matrices before and after adding an

• Then  $P' - P = \mathbf{x}^{\mathsf{T}_X} \mathbf{\delta}$ , where contains -1/d(d + 1) in positions of successors of x

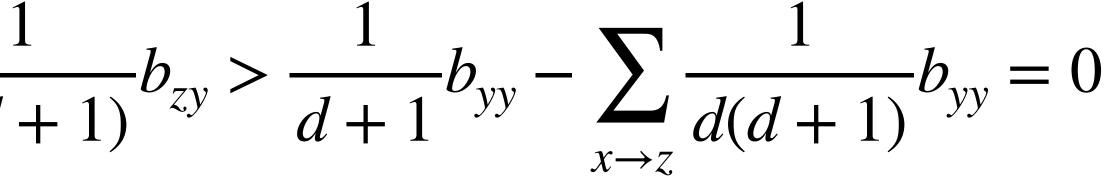
correction to  $(1 - \alpha P)^{-1}$ , but we can gather the ugly stuff in a positive constant

- McDonald, Neumann, Schneider & Tsatsomeros proved in 1995 that diagonal elements, so if  $B = (1 - \alpha P)^{-1}$ :

$$\left[\delta(1-\alpha P)^{-1}\right]_{y} = \frac{1}{d+1}b_{yy} - \sum_{x \to z}\frac{1}{d(d-x)}$$

 But then we can rewrite the difference between the PageRank vectors as  $v(1 - \alpha P')^{-1} - v(1 - \alpha P)^{-1} = (v + c\delta)(1 - \alpha P)^{-1} - (1 - \alpha P)^{-1} = c\delta(1 - \alpha P)^{-1}$ 

diagonal elements of the inverse of an M-matrix dominate strictly the off-



- A similar approach can be used to prove strict rank monotonicity
- There is no such proof in sight using Markov chains
- In general, properties that are not true on all Markov chains will not be within reach of proof techniques based on Markov chains
- Also, most strong results require regularity
- Linear algebra makes it possible to prove stronger statements using just irreducibility (strong connection)

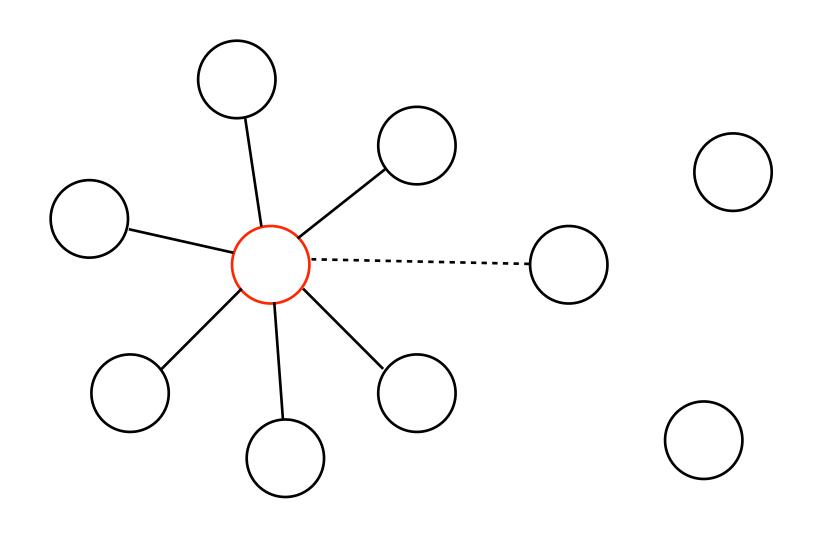
# Getting a New Friend

# **Seeley's: Score and Strictly Rank Monotone**

- Standardized Seeley's ( $\ell_1$ -normalized degree)
- Obviously strictly rank monotone
- Score monotone except in the case of a star plus isolated nodes

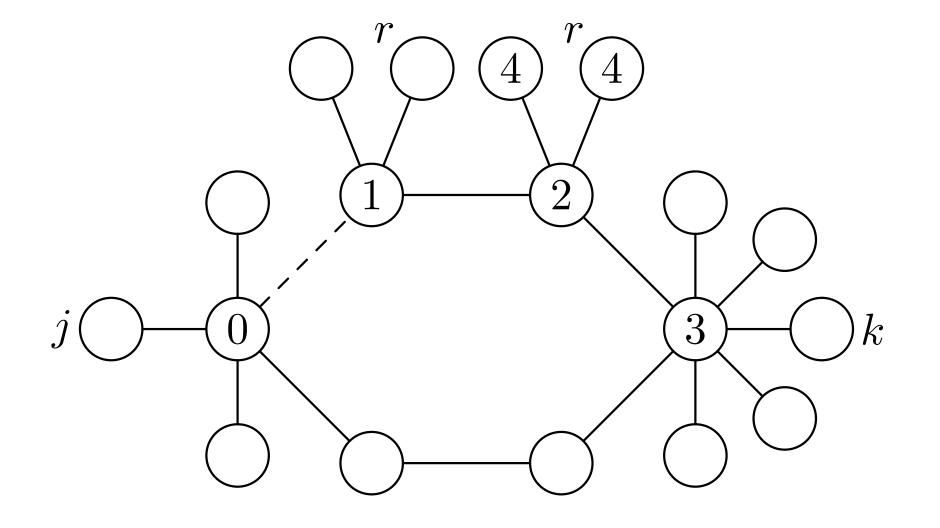
$$\frac{d(x) + 1}{2m + 2} > \frac{d(x)}{2m} \implies d(x)$$





# Nothing Else Works

• Geometric centralities



- nodes around 0

## • Idea: connecting 0 and 1 benefits 1 a lot, because it gets closer to the j

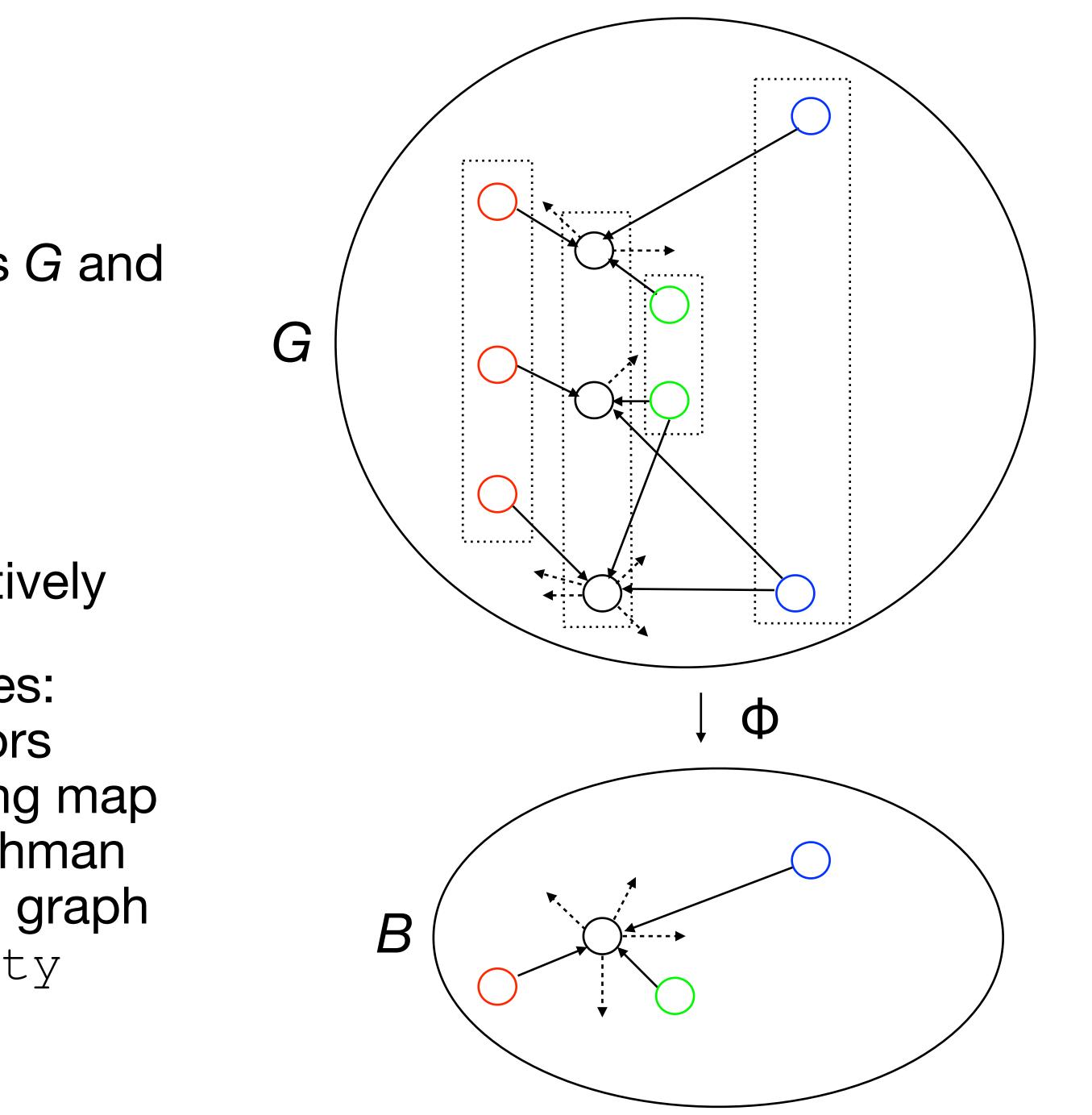
The same doesn't happen for 0, which gains very little as r is small

# PageRank: None of the Above, Again

- Much more complex situation
- We would like to have results for every value of the damping factor  $\alpha$ • For  $\alpha$  sufficiently close to 1, adding an edge never violates rank monotonicity
- Thus, different values of  $\alpha$  require different counterexamples
- This might suggest that in some right subinterval of the unit interval rank monotonicity works
- Surprisingly, this is not the case

# Fibrations of Graphs

- Maps (morphisms) between graphs G and B
- Fiber: counterimage of a node in B
- Local in-isomorphism property: inneighbourhoods are mapped bijectively
- Appeared under a plethora of names: equitable partitions and front divisors (spectral graph theory), left-resolving map (symbolic dynamics), Weisfeiler–Lehman test (graph isomorphism problems, graph neural networks), first stage of nauty (graph isomorphism), etc.



# **Fundamental Property**

- Take a fibration  $\phi: G \rightarrow B$
- If v is a vector on B, write  $v^{\phi}$  for the *lifting of* v along  $\phi$ , that is, the vector obtained by copying v along the fibers
- Then,  $(vB)^{\phi} = v^{\phi}G$  (Sachs 1966)
- the fiber values by B and lifting
- the theory of equitable partitions and graph divisors in the '60s)
- More importantly,  $(v \sum_{n \ge 0} \alpha^n B^n)^{\phi} = v^{\phi} \sum_{n \ge 0} \alpha^n G^n$  as lifting is linear



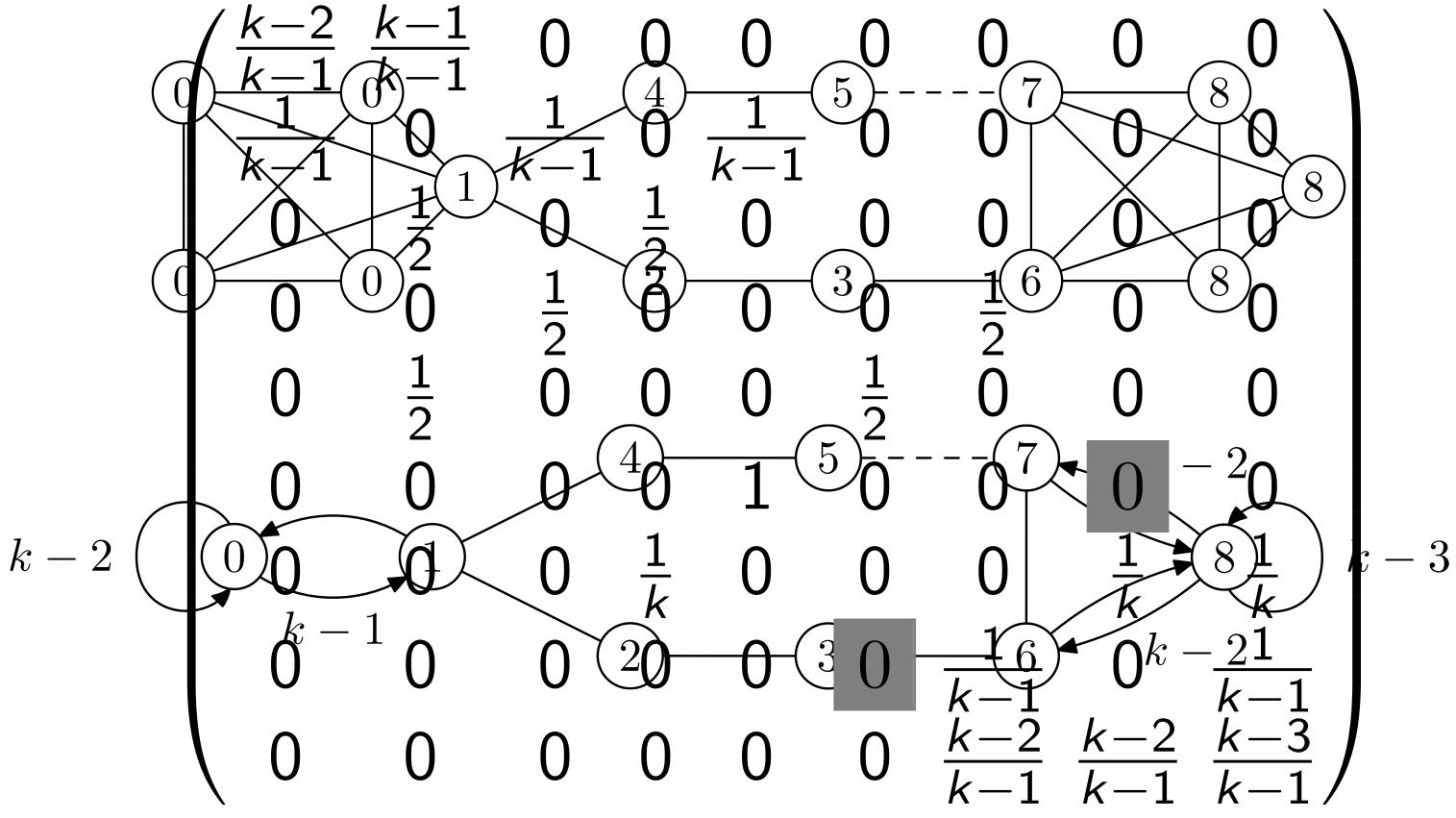
• That is: if a vector is fiberwise constant, you can multiply by G by multiplying

• Consequence: if G is strongly connected, the dominant eigenvector of G is the lifting of that of B, and the characteristic polynomial of B divides that of G (see



# Idea: Parametric Counterexample

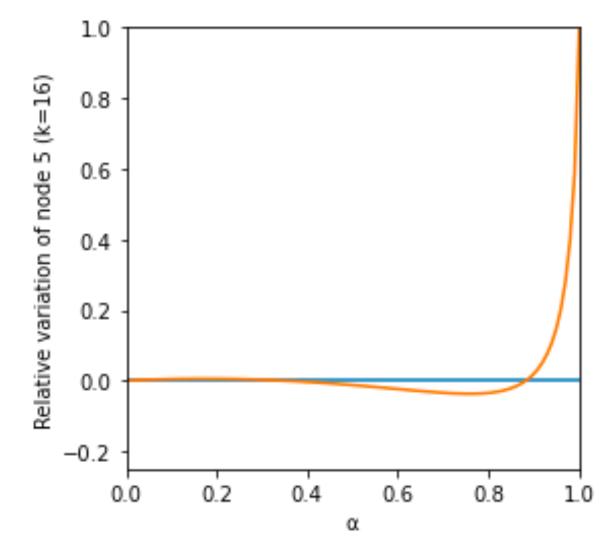
 We just need to find one that is fibe vertices



We just need to find one that is fibered over a graph with a finite number of

# Exact PageRank scores

- Now PageRank scores are rational functions in α whose coefficients are rational functions in k
- - Then we use Sturm polynomials to show that there are two sign changes around the middle
  - We sandwich the sign changes between points going to 0 and 1 for  $k \rightarrow \infty$
  - We get a counterexample for each value of  $\boldsymbol{\alpha}$

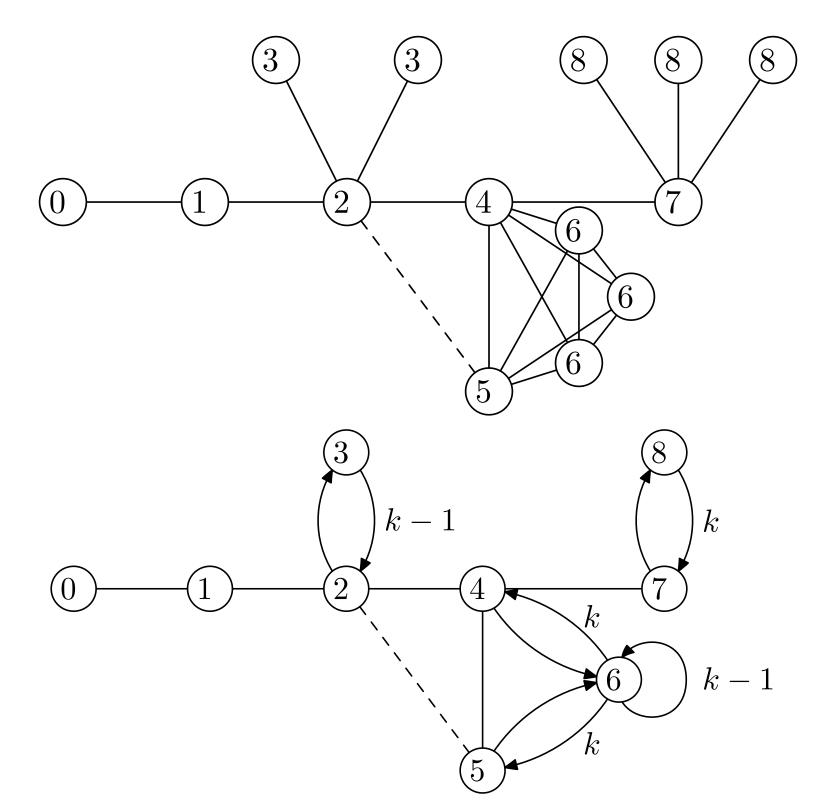




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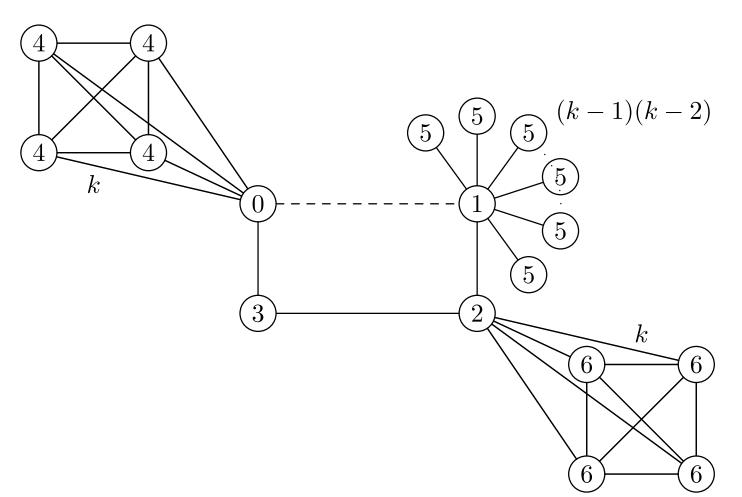
# **Rich Get Richer**

- In the previous example, the important node becomes more important • ("vampire" behavior)
- We have a similar class of counterexamples for the opposite behavior



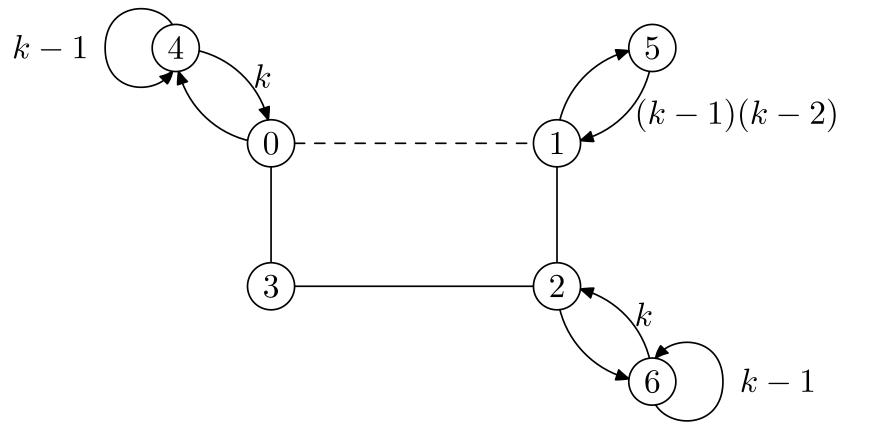
# **Eigenvector and Katz Centrality**

- We can use the same technique, with a different graph, to prove similar results for eigenvector centrality and Katz centrality



where p is the spectral radius.

Eigenvector centrality is particularly tricky because we cannot have an exact description of the dominant eigenvector (unless you have less than five nodes, or you're very lucky)



We solve the problem by proving the results for Katz's index in an interval around  $1/\rho$ ,

# It Actually Happens (IMDB)!

| Score increase    | Score decrease    | Rank violations                         |
|-------------------|-------------------|---|
| Meryl Streep      | Yasuhiro Tsushima | Anne–Mary Brown,<br>Jill Corso,         |
| Denzel Washington | Corrie Glass      | Patrice Fombelle,<br>John Neiderhauser, |
| Sharon Stone      | Mary Margaret (V) | Dolores Edwards, Colette<br>Hamilton,   |
| John Newcomb      | Robert Kirkham    | Brandon Matsui, Evis<br>Trebicka,       |

# Conclusion

- on directed graph is correct
- the directed case: Sherman–Morrison formula plus M-matrices
- graphs
- <u>http://vigna.di.unimi.it/fibrations/</u>

• The natural insight into the behavior of popular spectral centrality measures

General proof technique for monotonicity properties of spectral measures in

The same insight or proof techniques does not work anymore for undirected

General proof technique for spectral counterexamples: graph fibrations